

# 1.3 Algorithms and Convergence

# Algorithm & Pseudocode

- An **algorithm** is an ordered sequence of unambiguous and well-defined instructions that performs some tasks.
- **Pseudocode** is an **artificial and informal high-level language** that describes the operating principle of a computer program or algorithm.
  - Pseudocode allows ones to focus on the logic of the algorithm without being distracted by details of language syntax.
  - The pseudo-code is a "text-based" detail (algorithmic) design tool and is complete. It describes the entire logic of the algorithm so that implementation is a task of translating line by line into source code.
  - Pseudocode also uses structured programming design.

- Rules of pseudocode

1. Three categories of algorithmic operations

- a) sequential operations (Sequence) - instructions are executed in order.

- Example: "variable" = "expression".

- b) conditional operations (If-Then-Else) - a control structure that asks a true/false question and then selects the next instruction based on the answer.

- Example:

```
if "condition" then
```

```
    (subordinate) statement 1
```

```
else
```

```
    (subordinate) statement 2
```

- c) iterative (loop) operations (While) - a control structure that repeats the execution of a block of instructions

- Example:

```
while "condition"
```

```
    (subordinate) statement 1
```

```
    (subordinate) statement 2
```

2. All statements showing "dependency" are to be indented.

3. A period (.) indicates the termination of a step.

4. A semicolon (;) separates tasks within a step.

# Pseudocode Structure

INPUT:

OUTPUT:

*Step1:*

*Step2:*

*etc...*

# Pseudocode

**Example.** Compute  $\sum_{i=1}^N x_i$

INPUT  $N, x_1, x_2, \dots, x_N.$

OUTPUT  $SUM = \sum_{i=1}^N x_i$

*Step 1* **Set**  $SUM = 0.$  // Initialize accumulator

*Step 2* **For**  $i = 1, 2, \dots N$  **do**

    set  $SUM = SUM + x_i.$  // add next term

*Step 3* OUTPUT(SUM);

STOP.

# Characterizing Algorithms

## Error Growth

Suppose  $E_0 > 0$  denotes an initial error, and  $E_n$  is the error after  $n$  subsequent operations.

1. If  $E_n \approx CnE_0$ , where  $C$  is a const. independent of  $n$ : the growth of error is **linear**.
2. If  $E_n \approx C^n E_0$ , where  $C > 1$ : the growth of error is **exponential**.

**Remark:** linear growth is unavoidable; exponential growth must be avoided.

## Stability

- Stable algorithm: small changes in the initial data produce small changes in the final result
- Unstable or conditionally stable algorithm: small changes in all or some initial data produce large errors

**Example a.** For any  $c_1$  and  $c_2$ ,  $p_n = c_1 \left(\frac{1}{3}\right)^n + c_2 3^n$  is the solution to the recursive equation

$$p_n = \frac{10}{3} p_{n-1} - p_{n-2}, \quad \text{for } n = 2, 3, \dots$$

Suppose  $p_0 = 1$  and  $p_1 = \frac{1}{3}$ . Use 5-digit rounding arithmetic to compute  $\{p_n\}$ . Is the procedure stable?

## Definition 1.18 Rate of convergence for sequences

Suppose  $\{\beta_n\}_{n=1}^{\infty}$  is a sequence converging to 0, and  $\{\alpha_n\}_{n=1}^{\infty}$  converges to a number  $\alpha$ . If a positive constant  $K$  exists with

$$|\alpha_n - \alpha| \leq K|\beta_n|, \quad \text{for large } n,$$

then  $\{\alpha_n\}_{n=1}^{\infty}$  is said to converge to  $\alpha$  with rate of convergence  $O(\beta_n)$ , indicated by  $\alpha_n = \alpha + O(\beta_n)$ .

Typical  $\{\beta_n\}_{n=1}^{\infty}$ :

$$\beta_n = \frac{1}{n^p} \quad \text{for some } p > 0$$

**Example 2.** Suppose that, for  $n \geq 1$ ,  $\alpha_n = \frac{n+1}{n^2}$  and

$\hat{\alpha}_n = \frac{n+3}{n^3}$ . Determine rates of convergence for these two sequences.

## Definition 1.19 Rate of convergence for functions

Suppose that  $\lim_{h \rightarrow 0} G(h) = 0$  and  $\lim_{h \rightarrow 0} F(h) = L$ .

If a positive constant  $K$  exists with

$$|F(h) - L| \leq K|G(h)|, \quad \text{for sufficiently small } h,$$

then  $F(h) = L + O(G(h))$ .

Typical  $G(h)$ :

$$G(h) = h^p \quad \text{for some } p > 0$$

**Example 3.** Use the third Taylor polynomial about  $h = 0$  to show that  $\cosh h + \frac{1}{2}h^2 = 1 + \mathcal{O}(h^4)$ .