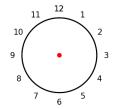
A Taste of Pi: Clocks, Set, and the Secret Math of Spies

Katherine E. Stange SFU / PIMS-UBC

October 16, 2010

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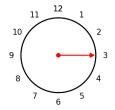
Here is a picture of a clock.



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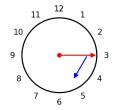




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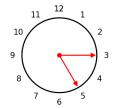


 $3\,pm + 2\,hours =$

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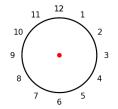


 $3\,pm + 2\,hours = 5\,pm$

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Here is a picture of a clock.

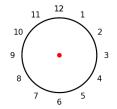


3 pm + 2 hours = 5 pm $3+2 \equiv 5 mod 12$

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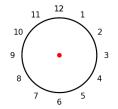
Here is a picture of a clock.



3 pm + 2 hours = 5 pm $3+2 \equiv 5 mod 12$ 2 pm + 11 hours =

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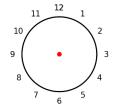


3pm+2hours	=	5 pm
3+2	\equiv	5 mod 12
2pm+11hours	=	1 am

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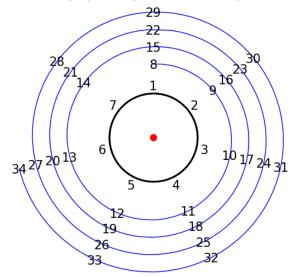
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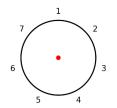
It's a little like rolling up a long line of the integers into a circle:



We could have a clock with any number of hours on it.

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Here is a picture of a clock with 7 hours.

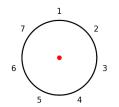


We could have a clock with any number of hours on it.

Here is a picture of a clock with 7 hours.

 $2 \, o' clock + 11 \, hours =$

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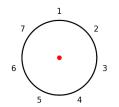


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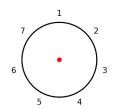
 $2 \, o' clock + 11 \, hours = 6 \, o' clock$

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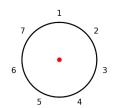


- $2 \, o' clock + 11 \, hours = 6 \, o' clock$
 - $2+11 \ \equiv \ 6 \ mod \ 7$

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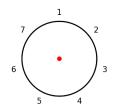
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$$1 \text{ o'clock} - 24 \text{ hours} =$$

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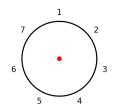
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1 o'clock - 24 hours = 5 o'clock

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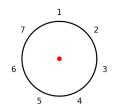


- $2 \, o' clock + 11 \, hours = 6 \, o' clock$
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- 1 o'clock 24 hours = 5 o'clock
 - $1-24 \ \equiv \ 5 \ mod \ 7$

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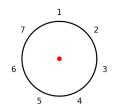
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 $2 \, \text{o'clock} \times 4 =$

We could have a clock with any number of hours on it.

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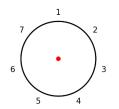


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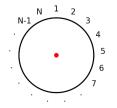
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 - $2 \times 4 \equiv 1 \mod 7$

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We could label these with days of the week ...



We call the *N*-hour clock \mathbb{Z}_N , and it has *N* elements:

$$\mathbb{Z}_N = \{0, 1, 2, 3, \dots, N-1\}$$

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We can add, subtract and multiply elements of \mathbb{Z}_N (and get back elements of \mathbb{Z}_N).

- ► The math of clocks is called *Modular Arithmetic* and *N* is called the *modulus*.
- Two numbers A and B are the same "modulo N" if A and B differ by adding N some number of times.

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We could say that a hamburger and a cheeseburger are the same modulo cheese.

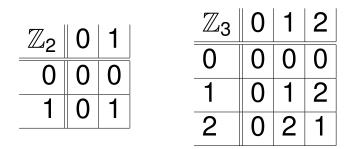
- The math of clocks is called *Modular Arithmetic* and *N* is called the *modulus*.
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- We could say that a hamburger and a cheeseburger are the same modulo cheese.
- Some people say Gauss invented modular arithmetic, but humans have used it as long as we've had...

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- clocks
- weeks
- gears
- money
- <u>ا ...</u>
- It's the beginning of the study of Number Theory.

Let the festivities begin!



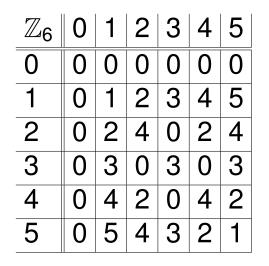


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\mathbb{Z}_4	0	1	2	3
0	0	0	0	0
1	0	1	2	3
2	0	2	0	2
3	0	3	2	1

\mathbb{Z}_5	0	1	2	3	4
0	0	0	0	0	0
1	0	1	2	3	4
2	0	2	4	1	3
3	0	3	1	4	2
4	0	4	3	2	1

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\mathbb{Z}_7	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0			3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2		1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

Z 8	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7
2	0	2	4	6	0	2	4	6
3	0	3	6	1	4	7	2	5
4	0	4	0	4	0	4	0	4
5	0	5	2	7	4	1	6	3
6	0	6	4	2	0	6	4	2
7	0	7	6	5	4	3	2	1

Th	e Matl	n of	Clo	cks	- M	ultip	lica	tion	Tab	les
	\mathbb{Z}_9	0	1		3	4	5	-	7	8
	0	0	0	0	0	0	0	0	0	0
	1	0	1	2	3	4	5	6	7	8
	2	0	2	4	6	8	1	3	5	7
	3	0	3	6	0	3	6	0	3	6
	4	0	4	8	3	7	2	6	1	5
	5	0	5	1	6	2	7	3	8	4
	6	0	6	3	0	6	3	0	6	3
	7	0	7	5	3	1	8	6	4	2
	8	0	8	7	6	5	4	3	2	

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\mathbb{Z}_{11}	0	1	2	3	4	5	6	7	8	9	10
0	0	0	0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7	8	9	10
2	0	2	4	6	8	10	1	3	5	7	9
3	0	3	6	9	1	4	7	10	2	5	8
4	0	4	8	1	5	9	2	6	10	3	7
5	0	5	10	4	9	3	8	2	7	1	6
6	0	6	1	7	2	8	3	9	4	10	5
7	0	7	3	10	6	2	9	5	1	8	4
8	0	8	5	2	10	7	4	1	9	6	3
9	0	9	7	5	3	1	10	8	6	4	2
10	0	10	9	8	7	6	5	4	3	2	1

1. When *N* is a prime number, then you can divide in \mathbb{Z}_N .

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- 2. This makes \mathbb{Z}_N a really great number system: it has $+, -, \times$ and \div .

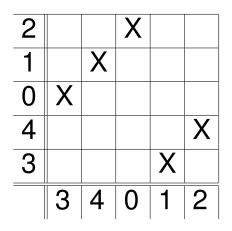
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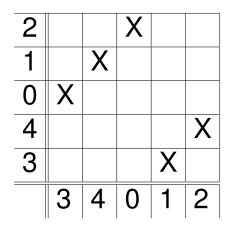
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3. It's even better than the integers (there's no 1/2 in the integers!).

The graph of the line y = x + 2 in \mathbb{Z}_5 :

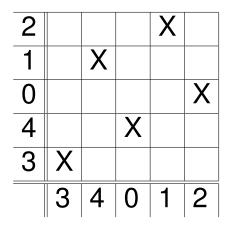


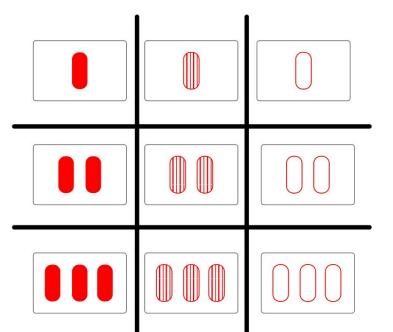
The graph of the line y = x + 2 in \mathbb{Z}_5 :



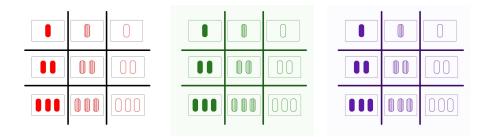
The graph is a little like Asteroids!

The graph of the line y = 3x + 4 in \mathbb{Z}_5 :





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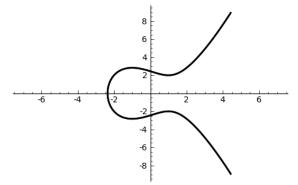
Set images due to Diane Maclagan and Ben Davis

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http://www.setgame.com/

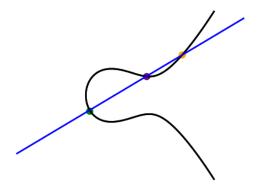


Here's the graph of $y^2 = x^3 - 3x + 6$ in the usual world (real numbers):

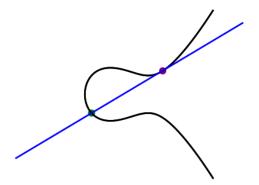


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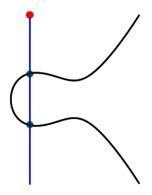
Adding two points to get another: P + Q + R = O.



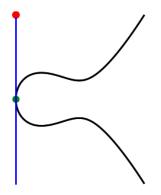
Adding a point and its negative: P + Q + Q = O.



Adding a point and its negative: P + -P = O.



A point which adds with itself to zero: P + P = O.



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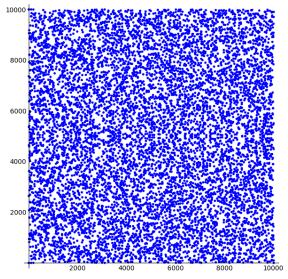
The graph of $y^2 = x^3 + 2x + 1$ in \mathbb{Z}_5 :

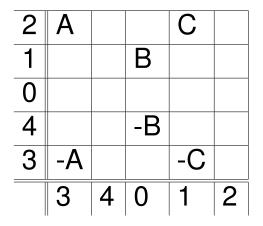
2	X			Х	
1			Х		
0					
4			Х		
3	Χ			Χ	
	3	4	0	1	2

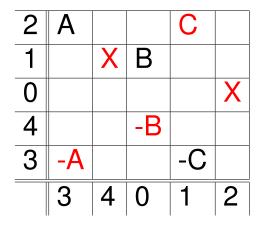
This is called an "Elliptic Curve"

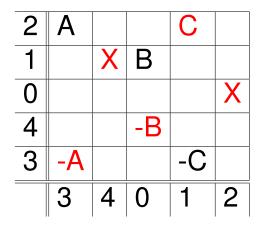
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Here's an elliptic curve in \mathbb{Z}_{10007} .

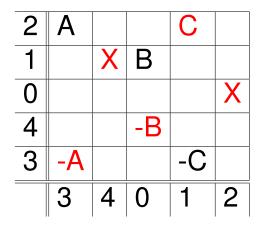




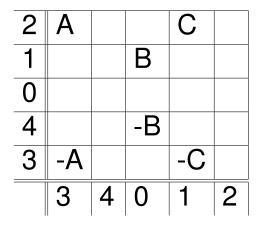


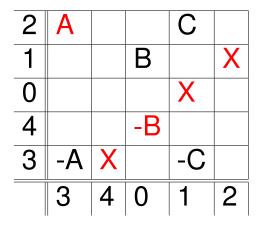


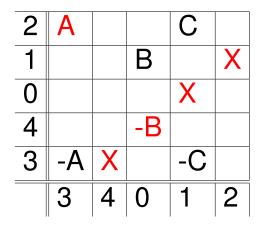
-A + -B + C = 0



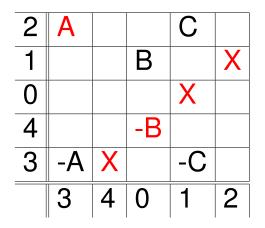
$$-A + -B + C = 0$$
$$A + B = C$$



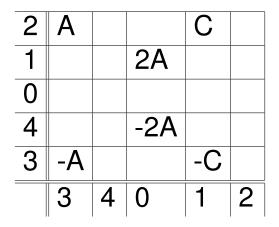




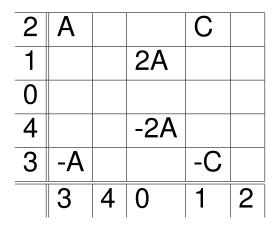
A + A + -B = 0



$$A + A + -B = 0$$
$$A + A = B$$



$$A + A + -B = 0$$
$$A + A = B$$
So
$$B = 2A$$

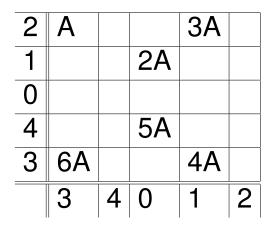


A + A + -B = 0 A + A = BSo B = 2AFrom last slide: C = A + B =A + 2A = 3A

2	A			ЗA	
1			2A		
0					
4			-2A		
3	-A			-3A	
	3	4	0	1	2

A + A + -B = 0 A + A = BSo B = 2AFrom last slide: C = A + B = A + 2A = 3ASo C = 3A

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With a little more work, we find out that -3A = 4A, -2A = 5A and -A = 6A, and finally that 7A = O.

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The Math of Spies - Elliptic Curve Addition Table

E	0	А	2A	3A	4A	5A	6A
0	Ο	Α	2A	ЗA	4A	5A	6A
Α	A	2A	ЗA	4A	5A	6A	0
2A	2A	ЗA	4A	5A	6A	0	Α
ЗA	3A	4A	5A	6A	0	Α	2A
4A	4A	5A	6A	0	А	2A	ЗA
	5A						
6A	6A	0	Α	2A	ЗA	4A	5A

The Math of Spies - Modular Arithmetic Addition Table

\mathbb{Z}_7	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

Suppose I gave you two numbers, P = 9 and Q = 45 and I said,

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Suppose I gave you two numbers, P = 9 and Q = 45 and I said,

"How many times need I add P to itself to get Q?"

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Suppose I gave you two numbers, P = 9 and Q = 45 and I said,

"How many times need I add P to itself to get Q?"

You would divide 45 by 9 and get the answer: 5. Division is fairly easy for integers!

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Suppose I gave you two numbers, P = 9 and Q = 45 and I said,

"How many times need I add P to itself to get Q?"

- You would divide 45 by 9 and get the answer: 5. Division is fairly easy for integers!
- It takes more time as the numbers get bigger, but the time it takes grows with the number of digits of the numbers.

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Suppose I gave you two numbers modulo 5, P = 2 and Q = 3 and I said,

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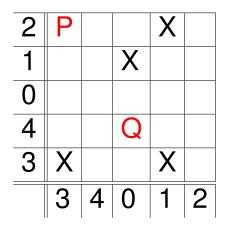
This is trickier. You could complete a multiplication table and look in it to search for the answer. It turns out there are faster ways.

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Suppose I gave you two numbers modulo 5, P = 2 and Q = 3 and I said,

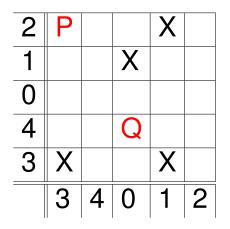
"How many times need I add P to itself to get Q?"

- This is trickier. You could complete a multiplication table and look in it to search for the answer. It turns out there are faster ways.
- The smartest algorithms (can you come up with one?), are about as fast as division for the integers. The time it takes grows with the number of digits of the modulus.



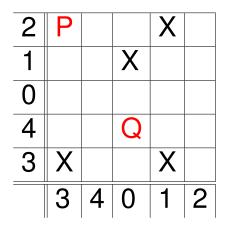
Suppose I gave you the points P and Q and I said "How many times need I add P to itself to get Q?"

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- Suppose I gave you the points P and Q and I said "How many times need I add P to itself to get Q?"
- ➤ You might remember that we found Q = 5P from our multiplication table.

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- Suppose I gave you the points P and Q and I said "How many times need I add P to itself to get Q?"
- You might remember that we found Q = 5P from our multiplication table.
- But it was a lot of work! Is there an easy way to do this?

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No one knows any efficient way to solve this problem !!



No one knows any efficient way to solve this problem!!

The time taken by good algorithms grows with about the square root of the size of the modulus.

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Modern cryptography is based on mathematical operations that are easy to do and hard to undo. Example:

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Getting pregnant.

Modern cryptography is based on mathematical operations that are easy to do and hard to undo. Example:

- Getting pregnant.
- Multiplying numbers is easy, but factoring them is hard.

Modern cryptography is based on mathematical operations that are easy to do and hard to undo.

Example:

- Getting pregnant.
- Multiplying numbers is easy, but factoring them is hard.
- On an elliptic curve, adding a point P to itself many times is easy. Figuring out how many times it was added (if you weren't watching) is hard. This is the elliptic curve discrete logarithm problem.

Alice and Bob want to share a secret.

Alice and Bob want to share a secret.

A point P on an elliptic curve is general knowledge.

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Alice Bob

Alice and Bob want to share a secret.

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Alice Bob secret a b

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	Alice	Bob
secret	а	b
public	аP	bP

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	Alice	Bob
secret	а	b
public	аP	bP

Alice and Bob can both compute *abP*.

Alice and Bob want to share a secret.

A point P on an elliptic curve is general knowledge.

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	Alice	Bob
secret	а	b
public	аP	bP

Alice and Bob can both compute *abP*.

No one else can compute it!

Here the size of the modulus N we use for this algorithm in your web browser, when you log onto a secure site:

N = 68647976601306097149819007990813932172694353

0014330540939446345918554318339765605212255964066 1454554977296311391480858037121987999716643812574 028291115057151

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The hard problem of factoring is used for cryptography called RSA.

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- The elliptic curve discrete logarithm problem is used for elliptic curve cryptography (ECC).

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- Together, these two hard problems are used for pretty nearly all the cryptography in the modern world: your bank, your cell phone, your computer.

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 (No one has come up with a security method based on pregnancy.)

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- If you can come up with a fast algorithm for these hard problems, you would immediately become hugely famous, you would get job offers from every government in the world, and would get invited on Oprah.

Thank you!

- Thanks to SFU, Veselin Jungic, Malgorzata Dubiel and Nadia Nosrati, and Jonathan Wise.
- And to you! Feel free to email me anytime (email on my website).

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